# Functional Programming in Scala Part II 

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## Tail Recursion

## Evaluating a Function Application

function call $f(e 1, \ldots, e n)$ is evaluated as follows:

- evaluate expressions e1,...,en resulting in values v1,...,vn
- replace $f(e 1, \ldots, e n)$ by body of function in which actual arguments replace formal parameters of $f$.

More formally,
$\operatorname{def} f(x 1, \ldots x n)=B$
$f(v 1, \ldots, v n)->[v 1 / x 1, \ldots, v n / x n] B$
where $[\mathrm{v} 1 / \mathrm{x} 1, \ldots, \mathrm{vn} / \mathrm{xn}] \mathrm{B}$ stands for expression B in which all occurrences of xi are replaced by vi.
$[\mathrm{v} 1 / \mathrm{x} 1, \ldots, \mathrm{vn} / \mathrm{xn}]$ is called a substitution.

## Rewriting Example (gcd)

def gcd(a: Int, b: Int): Int =
if $(b==0)$ a else $\operatorname{gcd}(b, a \% b)$
$\operatorname{gcd}(14,21)$
$\rightarrow>$ if $(21==0) 14$ else $\operatorname{gcd}(21,14 \% 21)$
$\rightarrow>$ if (false) 14 else $\operatorname{gcd}(21,14 \% 21)$
$\rightarrow>\operatorname{gcd}(21,14 \% 21)$
$\rightarrow>\operatorname{gcd}(21,14)$
$\rightarrow>$ if $(14==0) 21$ else $\operatorname{gcd}(14,21 \% 14)$
$\rightarrow>$ if (false) 21 else $\operatorname{gcd}(14,21 \% 14)$
$\rightarrow>\operatorname{gcd}(14,21 \% 14)$
$\rightarrow \operatorname{gcd}(14,7)$
$\rightarrow>$ if $(7==0) 14$ else $\operatorname{gcd}(7,14 \% 7)$
$\rightarrow>$ if (false) 14 else $\operatorname{gcd}(7,14 \% 7)$
$->\operatorname{gcd}(7,14 \% 7)$
$\rightarrow>\operatorname{gcd}(7,0)$
$\rightarrow>$ if $(0==0) 7$ else $\operatorname{gcd}(0,7 \% 0)$
$\rightarrow>7$

## Tail Recursion is better

If a function calls it self as its last action, the function's stack can be reused. This is called tail recursion

Tail recursive functions are essentially iterative processes
In Scala, tail-recursive functions can be annotated

```
@tailrec
def gcd(a: Int, b: Int): Int = ...
```

Tail recursive version of factorial:
def fact(answer: Int, n: Int): Int = if ( $\mathrm{n}==0$ ) answer else fact( $\mathrm{n}^{*}$ answer, $\mathrm{n}-1$ )
def factorial(n: $\operatorname{lnt}$ ): $\operatorname{Int}=$ fact(1,n)

## Higher Order Functions

Functional languages treat functions as first-class values.
i.e. like any other value, a function can be passed as a parameter and returned as a result.

This provides a flexible way to compose programs.
Functions that take other functions as parameters or return functions as results are called higher-order functions.

## Higher Order Functions - Example

```
def sumlnts(a: Int, b: Int): Int =
    if ( \(a>b\) ) 0 else \(a+\operatorname{sum} \operatorname{lnts}(a+1, b)\)
def cube( \(x\) : \(\ln\) ) : \(\operatorname{Int}=x^{*} x^{*} x\)
def fact( \(x: \operatorname{lnt}):\) Int \(=\) if \((x==0) 1\) else \(x\) * fact( \(x-1\) )
def sumCubes(a: Int, b : Int): Int =
    if \((a>b) 0\) else cube \((a)+\operatorname{sumCubes}(a+1, b)\)
def sumFactorials(a: Int, b: Int): Int =
    if \((a>b) 0\) else fact \((a)+\operatorname{sumFactorials}(a+1, b)\)
```

Can we factor out the function and reduce all of these to a single function?

## Higher Order Functions - Continued

```
def sum(f: Int => Int, a: Int, b: Int): Int =
    if (a>b) 0 else f(a) + sum(f, a+1,b)
def sumInts(a: Int, b: Int): Int = sum(id, a, b)
def sumCubes(a: Int, b: Int): Int = sum(cube, a, b)
def sumFactorials(a: Int, b: Int): Int = sum(fact, a, b)
where
def id(x: Int): Int = x
def cube(x: Int): Int = x * x * x
def fact(x: Int): Int = if (x == 0) 1 else x * fact(x - 1)
Type \(A=>B\) is the type of a function that takes an argument of type \(A\) and returns a result of type \(B\).
```

So, Int => Int is the type of functions that map integer to integers

## Anonymous Functions

Passing functions as parameters leads to the creation of many small functions.

It is tedious to have to define (using def) and name these functions.
e.g.
def str = "abc";
println(str);
vs
printIn("str")
Can we not have function literals just like String literals?
Anonymous functions are basically function literals

## Anonymous Function Syntax

## Examples:

( $\mathrm{x}: \ln \mathrm{t})=>\mathrm{x}^{*} \mathrm{x}$ * x
( x : Int) is the parameter of the function and x * x * x is the body.
( x : $\ln \mathrm{t}, \mathrm{y}: \operatorname{Int}$ ) $=>\mathrm{x}+\mathrm{y}$
In general:
( $\mathrm{x} 1: \mathrm{T} 1, \mathrm{x} 2: \mathrm{T} 2, \ldots, \mathrm{xn}: \mathrm{Tn}$ ) => E can be expressed using def as follows:
$\{\operatorname{def} f(x 1: T 1, x 2: T 2, \ldots, x n: T n)=E ; f\}$
Using anonymous functions, we can write earlier sums in a shorter way:
def sumlnts(a: Int, $\mathrm{b}: \operatorname{lnt})=\operatorname{sum}(\mathrm{x}=>\mathrm{x}, \mathrm{a}, \mathrm{b})$
def sumCubes(a: Int, $b$ : Int $)=\operatorname{sum}\left(x=>x^{*} x^{*} x, a, b\right)$

## Products, Factorials, etc

```
def sum(a: Int, b: Int): Int =
    if (a>b)0 else a + sum(a+1,b)
def product(a: Int, b: Int): Int =
    if (a>b) 1 else a * product(a+1,b)
```

def factorial(n: $\operatorname{lnt}):$ Int $=\operatorname{product}(1, \mathrm{n})$

```
object SumProduct {
    def operate(f: (Int, Int)=>Int, ident: Int, a: Int, b: Int): Int =
        if (a > b) ident else f(a, operate(f, ident, a+1,b))
    def sum(a: Int, b: Int): Int =
        operate((x,y)=>x+y, 0, a, b)
    def product(a: Int, b: Int): Int =
        operate((x, y)=>x*y, 1, a, b)
    def main(args: Array[String]) {
        println(sum(1, 6))
        println(product(1, 6))
    }
}
```


## Currying

## Motivation:

def sumlnts(a: Int, b: Int): Int $=\operatorname{sum}(x=>x, a, b)$
def sumCubes(a: Int, b: Int): Int = sum( $\left.x=>x^{*} x^{*} x, a, b\right)$
def sumFactorials(a: Int, b: Int): Int = sum(fact, a, b)
parameters $a$ and $b$ get passed on to sum() without any modifications. Can we get rid of these parameters?

Function returning functions:

```
def sum(f: Int => Int): (Int, Int) => Int = {
    def sumF(a: Int, b: Int): Int =
        if (a>b)0 else f(a) + sumF(a+1,b)
    sumF
}
```

sum is a function that returns another function, sumF. sumF applies the given function parameter $f$ and sums the results.

## Currying continued

With the new definition of sum, we can define

```
def sumlnts = sum(x => x)
def sumCubes = sum(x => x***x)
def sumFactorials = sum(fact)
```

and use them as follows:
sumCubes $(1,10)+$ sumFactorials $(1,5)$
We can even avoid the middlemen sumlnt, sumCubes etc...
sum(cube) $(1,10)$
sum(cube) returns the sum of cubes function and this function is next applied to arguments $(1,10)$.

Function applications associate to the left:
sum(cube) $(1,10)=($ sum $($ cube $))(1,10)$

## Multiple Parameter Lists - Scala Syntax

Special Syntax in Scala (the following is equivalent to the nested sumF):
def sum(f: Int => Int)(a: Int, b: Int): Int = if $(a>b) 0$ else $f(a)+\operatorname{sum}(f)(a+1, b)$

In general
$\operatorname{def} f(\operatorname{args} 1) \ldots(\operatorname{argsn})=E, n>1$ is equivalent to
$\operatorname{def} f(\operatorname{args} 1) \ldots(\operatorname{argsn}-1)=\{\operatorname{def} g(\operatorname{argsn})=E ; g\}$
where g is a new function symbol or even shorter:
$\operatorname{def} f(\operatorname{args} 1) \ldots(\operatorname{argsn}-1)=(\operatorname{argsn}=>E)$
Repeating this n times, we get
$\operatorname{def} \mathrm{f}=(\operatorname{args} 1=>(\operatorname{args} 2=>\ldots(\operatorname{argsn}=>E) . .)$.
"Currying"

## Example

Given
def sum(f: Int => $\operatorname{Int})(\mathrm{a}: \operatorname{Int}, \mathrm{b}: \operatorname{Int}): \operatorname{Int}=\ldots$
what is the type of sum?

Answer:
(Int => Int) => (Int, Int) => Int
Since function types associate to the right, this can be rewritten as
$(\operatorname{lnt}=>\operatorname{lnt})=>((\operatorname{lnt}, \operatorname{lnt})=>\operatorname{lnt})$

## Problem - Higher Order Function

Write a product function similar to sum.
Generalize sum and product using HOFs

```
object SumProduct {
    def operate(f: (Int, Int)=> Int, ident: Int, a: Int, b: Int): Int =
        if (a > b) ident else f(a,operate(f,ident,a+1,b))
    def sum(a: Int, b: Int): Int =
        operate((x,y)=>x+y, 0, a, b)
    def product(a: Int, b: Int): Int =
        operate((x,y)=>x*y, 1, a, b)
    def main(args: Array[String]) {
        println(sum(1, 6))
        println(product(1, 6))
    }
}
```


## Example: Finding Fixed Points

A number $x$ is called a fixed point of a function $f$ if
$f(x)=x$
Very useful concept in computer science! For some functions, we can locate the fixed point by starting with an initial estimate, and then applying $f$ in a repetitive manner:
$x, f(x), f(f(x)), f(f(f(x))), \ldots$
until the value does not vary any more (or the change is sufficiently small)

## Scala Program to compute Fixed Points

```
val tolerance = 0.0001
def isCloseEnough(x: Double, y: Double) =
    abs((x-y)/x) < tolerance
def fixedPoint(f: Double => Double)(firstGuess: Double) = {
    def iterate(guess: Double): Double = {
        val next = f(guess)
        println(next)
        if (isCloseEnough(guess, next)) next
        else iterate(next)
    }
    iterate(firstGuess)
}
sqrt(x) can be expressed in terms of a fixed point as follows:
sqrt(x) = the number y such that y * y = x
    = the number y such that }\textrm{y}=\textrm{x}/\textrm{y
    = fixed point of function ( }y=>x/y
def sqrt(x: Double) = fixedPoint(y => x/y)(1.0)
```


## Example: Finding Fixed Points

The previous solution goes into an infinite loop!
Values oscillate between 2 and 1!
To fix this, use the function $(y=>(y+x / y) / 2)$ which takes the average of guess and next guess.
def sqrt( $x$ : Double $)=$ fixedPoint $(y=>(y+x / y) / 2)(1.0)$
This produces:
$\operatorname{sqrt}(2)=1.4142135623746899$

## Functions as Return Values

This example illustrates return vales as functions:
Recall sqrt( $x$ ) is a fixed point of function ( $y=>x / y$ )
The iterative algorithm converges to a solution by averaging successive values. This technique of stabilizing by averaging can be generalized into a function!
def averageDamp(f: Double => Double)(x: Double) $=(x+f(x)) / 2$
def sqrt(x: Double) $=$ fixedPoint(averageDamp $(y=>x / y))(1.0)$
This expresses the algorithm precisely!

