

3. Give the truth value of each of the following wffs in the interpretation where the domain consists of the states of the United States,  $Q(x, y)$  is "x is north of y,"  $P(x)$  is "x starts with the letter M," and  $a$  is "Massachusetts."
- $(\forall x)P(x)$
  - $(\forall x)(\forall y)(\forall z)[Q(x, y) \wedge Q(y, z) \rightarrow Q(x, z)]$
  - $(\exists y)(\exists x)Q(y, x)$
  - $(\forall x)(\exists y)[P(y) \wedge Q(x, y)]$
  - $(\exists y)Q(a, y)$
  - $(\exists x)[P(x) \wedge Q(x, a)]$
4. Give the truth value of each of the following wffs in the interpretation where the domain consists of people:  $M(x, y)$  is "x is the mother of y,"  $F(x)$  is "x is female,"  $M(x)$  is "m is male."
- $(\forall x)(\exists y)(M(y, x))$
  - $(\exists x)(\forall y)(M(x, y))$
  - $(\forall x)(\forall y)(M(x, y) \rightarrow M(y))$
  - $(\exists x)(\exists y)(M(x, y) \wedge M(y))$
  - $(\exists x)(\forall y)(M(x, y) \rightarrow F(y))$
5. For each wff, find an interpretation in which it is true and one in which it is false.
- $(\forall x)(([A(x) \vee B(x)] \wedge [A(x) \wedge B(x)])')$
  - $(\forall x)(\forall y)[P(x, y) \rightarrow P(y, x)]$
  - $(\forall x)[P(x) \rightarrow (\exists y)Q(x, y)]$
  - $(\exists x)[A(x) \wedge (\forall y)B(x, y)]$
  - $[(\forall x)A(x) \rightarrow (\forall x)B(x)] \rightarrow (\forall x)[A(x) \rightarrow B(x)]$
6. Identify the scope of each of the quantifiers in the following wffs and indicate any free variables.
- $(\forall x)[P(x) \rightarrow Q(y)]$
  - $(\exists x)[A(x) \wedge (\forall y)B(y)]$
  - $(\exists x)[(\forall y)P(x, y) \wedge Q(x, y)]$
  - $(\exists x)(\exists y)[A(x, y) \wedge B(y, z) \rightarrow A(a, z)]$
7. Which of the following are equivalent to the statement  
All circles are round.
- If it's round, it's a circle.
  - Roundness is a necessary property of circles.
  - Something that isn't round can't be a circle.
  - Some round things are circles.
8. Which of the following are equivalent to the statement  
Cats are smarter than dogs.
- Some cats are smarter than some dogs.
  - There is a cat that is smarter than all dogs.
  - All cats are smarter than all dogs.
  - Only cats are smarter than dogs.
  - All cats are smarter than any dog.
9. Using the predicate symbols shown and appropriate quantifiers, write each English language statement as a predicate wff. (The domain is the whole world.)
- $D(x)$  is "x is a day."  
 $S(x)$  is "x is sunny."  
 $R(x)$  is "x is rainy."  
 $M$  is "Monday."  
 $T$  is "Tuesday."

- \* a. All days are sunny.
  - \* b. Some days are not rainy.
  - \* c. Every day that is sunny is not rainy.
  - d. Some days are sunny and rainy.
  - e. No day is both sunny and rainy.
  - f. It is always a sunny day only if it is a rainy day.
  - g. No day is sunny.
  - h. Monday was sunny; therefore every day will be sunny.
  - i. It rained both Monday and Tuesday.
  - j. If some day is rainy, then every day will be sunny.
10. Using the predicate symbols shown and appropriate quantifiers, write each English language statement as a predicate wff. (The domain is the whole world.)

$B(x)$  is "x is a ball."

$R(x)$  is "x is round."

$S(x)$  is "x is a soccer ball."

- a. All balls are round.
  - b. Not all balls are soccer balls.
  - c. All soccer balls are round.
  - d. Some balls are not round.
  - e. Some balls are round but soccer balls are not.
  - f. Every round ball is a soccer ball.
  - g. Only soccer balls are round balls.
  - h. If soccer balls are round, then all balls are round.
11. Using the predicate symbols shown and appropriate quantifiers, write each English language statement as a predicate wff. (The domain is the whole world.)

$P(x)$  is "x is a person."

$T(x)$  is "x is a time."

$F(x, y)$  is "x is fooled at y."

- a. You can fool some of the people all of the time.
  - b. You can fool all of the people some of the time.
  - c. You can't fool all of the people all of the time.
12. Using the predicate symbols shown and appropriate quantifiers, write each English language statement as a predicate wff. (The domain is the whole world.)

$L(x)$ : x is a lion.

$R(x)$ : x roars.

$P(x)$ : x is a predator.

$Z(x)$ : x is a zebra.

$E(x, y)$ : x eats y.

- a. All lions are predators.
- b. Some lions roar.
- c. Only lions roar.
- d. Some lions eat all zebras.
- e. All lions eat all zebras.

13. Using the predicate symbols shown and appropriate quantifiers, write each English language statement as a predicate wff. (The domain is the whole world.)

$G(x)$ :  $x$  is a game.

$M(x)$ :  $x$  is a movie.

$F(x, y)$ :  $x$  is more fun than  $y$ .

- Any movie is more fun than any game.
  - No game is more fun than every movie.
  - Only games are more fun than movies.
  - All games are more fun than some movie.
14. Using the predicate symbols shown and appropriate quantifiers, write each English language statement as a predicate wff. (The domain is the whole world.)

$J(x)$  is " $x$  is a judge."

$C(x)$  is " $x$  is a chemist."

$L(x)$  is " $x$  is a lawyer."

$W(x)$  is " $x$  is a woman."

$A(x, y)$  is " $x$  admires  $y$ ."

- There are some women lawyers who are chemists.
  - ★ No woman is both a lawyer and a chemist.
  - ★ Some lawyers admire only judges.
  - All judges admire only judges.
  - Only judges admire judges.
  - All women lawyers admire some judge.
  - Some women admire no lawyer.
15. Using the predicate symbols shown and appropriate quantifiers, write each English language statement as a predicate wff. (The domain is the whole world.)

$C(x)$  is " $x$  is a Corvette."

$P(x)$  is " $x$  is a Porsche."

$F(x)$  is " $x$  is a Ferrari."

$S(x, y)$  is " $x$  is slower than  $y$ ."

- ★ Nothing is both a Corvette and a Ferrari.
  - ★ Some Porsches are slower than only Ferraris.
  - Only Corvettes are slower than Porsches.
  - All Ferraris are slower than some Corvette.
  - Some Porsches are slower than no Corvette.
  - If there is a Corvette that is slower than a Ferrari, then all Corvettes are slower than all Ferraris.
16. Using the predicate symbols shown and appropriate quantifiers, write each English language statement as a predicate wff. (The domain is the whole world.)

19. Three forms of negation are given for each statement. Which is correct?
- ★ a. Some people like mathematics.
    - 1. Some people dislike mathematics.
    - 2. Everybody dislikes mathematics.
    - 3. Everybody likes mathematics.
  - b. Everyone loves ice cream.
    - 1. No one loves ice cream.
    - 2. Everyone dislikes ice cream.
    - 3. Someone doesn't love ice cream.
  - c. All people are tall and thin.
    - 1. Someone is short and fat.
    - 2. No one is tall and thin.
    - 3. Someone is short or fat.
  - d. Some pictures are old or faded.
    - 1. Every picture is neither old nor faded.
    - 2. Some pictures are not old or faded.
    - 3. All pictures are not old or not faded.
20. Three forms of negation are given for each statement. Which is correct?
- a. Nobody is perfect.
    - 1. Everyone is imperfect.
    - 2. Everyone is perfect.
    - 3. Someone is perfect.
  - b. All swimmers are tall.
    - 1. Some swimmer is not tall.
    - 2. There are no tall swimmers.
    - 3. Every swimmer is short.
  - c. Every planet is cold and lifeless.
    - 1. No planet is cold and lifeless.
    - 2. Some planet is not cold and not lifeless.
    - 3. Some planet is not cold or not lifeless.
  - d. No bears are hungry.
    - 1. Only bears are hungry.
    - 2. All bears are hungry.
    - 3. There is a hungry bear.
21. Write the negation of each of the following.
- a. Some Web sites feature audio.
  - b. Every Web site has both audio and video.
  - ★ c. Every Web site has either audio or video.
  - d. Some Web sites have neither audio nor video.
  - e. Every Web site either has text or else has both audio and video.
22. Write the negation of each of the following.
- a. Only students eat pizza.
  - b. Every student eats pizza.
  - c. Some students eat only pizza.
23. Write the negation of each of the following.
- a. Some farmer grows only corn.
  - b. All farmers grow corn.
  - c. Corn is grown only by farmers.

24. Explain why each wff is valid.
- $(\forall x)(\forall y)A(x, y) \leftrightarrow (\forall y)(\forall x)A(x, y)$
  - $(\exists x)(\exists y)A(x, y) \leftrightarrow (\exists y)(\exists x)A(x, y)$
  - $(\exists x)(\forall y)P(x, y) \rightarrow (\forall y)(\exists x)P(x, y)$
  - $A(a) \rightarrow (\exists x)A(x)$
  - $(\forall x)[A(x) \rightarrow B(x)] \rightarrow [(\forall x)A(x) \rightarrow (\forall x)B(x)]$
25. Give interpretations to prove that each of the following wffs is not valid:
- $(\exists x)A(x) \wedge (\exists x)B(x) \rightarrow (\exists x)[A(x) \wedge B(x)]$
  - $(\forall x)(\exists y)P(x, y) \rightarrow (\exists x)(\forall y)P(x, y)$
  - $(\forall x)[P(x) \rightarrow Q(x)] \rightarrow [(\exists x)P(x) \rightarrow (\forall x)Q(x)]$
  - $(\forall x)[A(x)]' \leftrightarrow [(\forall x)A(x)]'$
26. Decide whether each of the following wffs is valid or invalid. Justify your answer.
- $(\exists x)A(x) \leftrightarrow ((\forall x)[A(x)]')'$
  - $(\forall x)P(x) \vee (\exists x)Q(x) \rightarrow (\forall x)[P(x) \vee Q(x)]$
  - $(\forall x)A(x) \rightarrow ((\exists x)[A(x)]')'$
  - $(\forall x)[P(x) \vee Q(x)] \rightarrow (\forall x)P(x) \vee (\exists y)Q(y)$

## SECTION 1.4 Predicate Logic

We can imagine arguments of the form

$$P_1 \wedge P_2 \wedge P_3 \wedge \dots \wedge P_n \rightarrow Q$$

where the wffs are built from predicates and quantifiers as well as logical connectives and grouping symbols. For a **valid argument**,  $Q$  must follow logically from  $P_1, \dots, P_n$  based solely on the internal structure of the argument, not on the truth or falsity of  $Q$  in any particular interpretation. In other words, the wff

$$P_1 \wedge P_2 \wedge P_3 \wedge \dots \wedge P_n \rightarrow Q$$

must be valid—true in all possible interpretations. No equivalent of the truth table exists to easily prove validity, so we turn to a formal logic system called **predicate logic**. We again use a system of derivation rules to build a proof sequence leading from the hypotheses to the conclusion. The rules should once more be truth-preserving so that if in some interpretation all the hypotheses are true, then the conclusion will also be true in that interpretation. The system will then be correct (*only* valid arguments will be provable). We also want the system to be complete (*every* valid argument should be provable), yet at the same time the rule set should be minimal.

### Derivation Rules for Predicate Logic

The equivalence rules and inference rules of propositional logic are still part of predicate logic. An argument of the form

$$P \wedge (P \rightarrow Q) \rightarrow Q$$

is still valid by modus ponens, even if the wffs involved are predicate wffs.